

<codedaily/>

Data Structure
& Algorithms 101
with doodles



girliemac



@girlie_mac



$O(n)$



$O(n^2)$



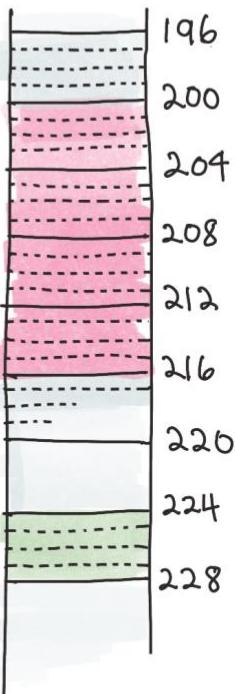
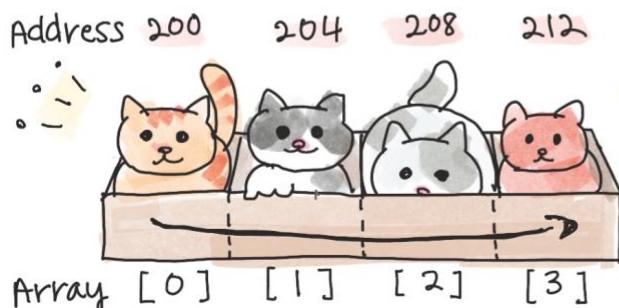
$O(\log n)$

Data Structures

Array & Linked List

Arrays

a linear data structure, stored in contiguous memory locations.



- (♥ Assume each 🐱 is an integer)
- = requires 4 bytes space
- ♥ The array of 🐱 must be allocated contiguously!
- address 200 — 216



- ♥ can randomly access w/ index
- $a[2] \rightarrow$ 🐱
- ♥ contiguous = no extra memory allocated = no memory overflow



- ☺ fixed size. Large space may not be avail for big array
- := 🐱 took the space! :=
- ☺ Insert & delete elements are costly.
 - may need to create a new copy of the array & allocate at a new address.

Big O Notation

Time complexity of an algorithm

"How much time it takes to run a function as
the size of the input grows."
Runtime

Const

array1 = [, , , , ]

array
number of elements
 $n=5$

Let's see if there is a needle in the haystack!

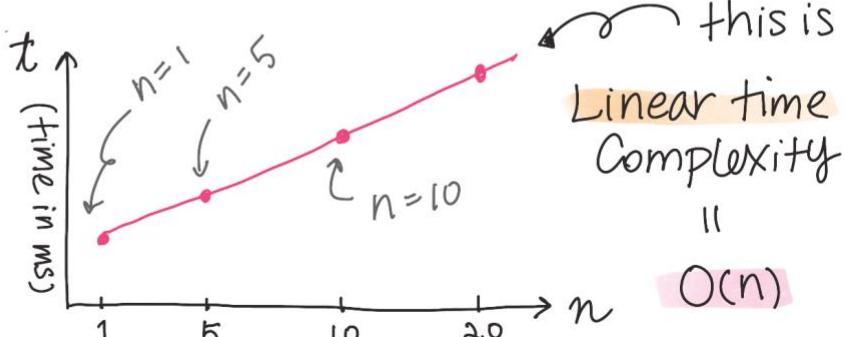
JS

```
Const numNeedles=(haystack, needle) => {
    let count=0
    for(let i=0; haystack.length; i++) {
        if(haystack[i] === needle) Count +=1;
    }
    return count;
```



How long does it take to execute when
the number of elements (n) is:

execution
time grows
linearly as
array size
increases!



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Big O Notation

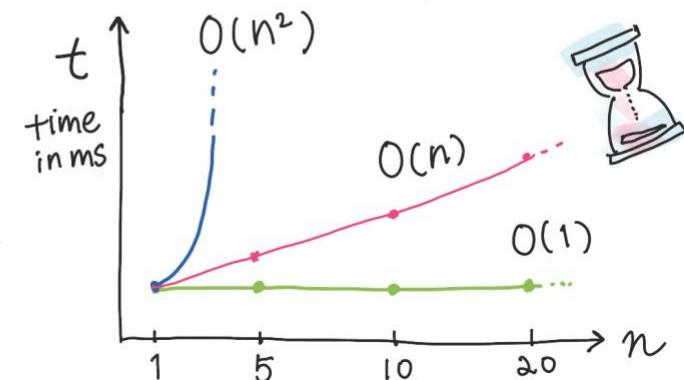
JS Let's see if we have some function that doesn't actually loop the array:

```
const alwaysTrueNoMatterWhat = (haystack) => {  
    return true;
```

$n=5$
 $n=10$
 $n=20$
⋮ Array size has no effect on the runtime

Constant time

||
 $O(1)$



Quadratic time = $O(n^2)$

Const

```
array2 = [ , , , ,  ];
```

$n=5$, however the runtime is proportional to n^2

JS Const hasDuplicates = (arr) => {
 for (let i = 0; i < arr.length; i++)
 let item = arr[i];
 if (arr.slice(i+1).indexOf(item) != -1) {
 return true;
 }
 return false;
}

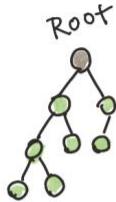
① Loop thru the array
② Another array lookup w/ indexOf method

Data Structures

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BST

Binary Search Tree



Binary tree

- tree data structure
- each node has at most 2 children

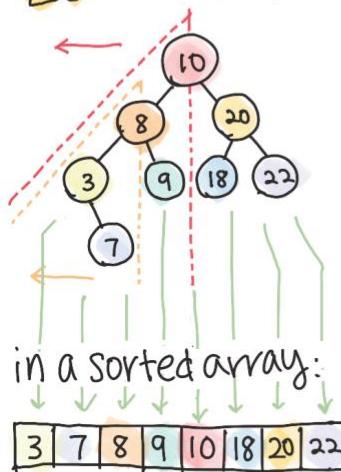
Binary heap

Binary Search Tree

a.k.a. Ordered or sorted binary tree

fast lookup
e.g. phone number lookup table by name

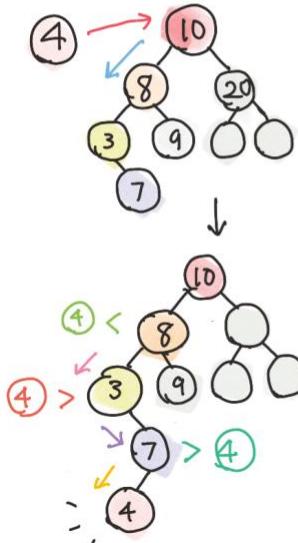
Rule of thumb



- each value of all nodes in the left subtrees is lesser
- \triangle 10's left subtrees: 8, 3, 9, 7
- \triangle 8: 3, 7 ← smaller than parent
- each value of all nodes in the right subtrees is larger
- no duplicate values

Insertion

→ Always add to the lowest spot to be a leaf No rearrange!



Let's add 4

1. Compare w/ the root first.
 2. $4 < 10$ so go left.
 3. then compare w/ the next, 8
 4. $4 < 8$ so go left
 5. Compare w/ the 3
 6. $4 > 3$ so go right.
 7. Compare w/ the 7
 8. $4 < 7$, so add to the left! Done.
- Complexity:
Ave. $O(\log n)$
Worst. $O(n)$

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Data Structures

BST

Binary Search Tree!

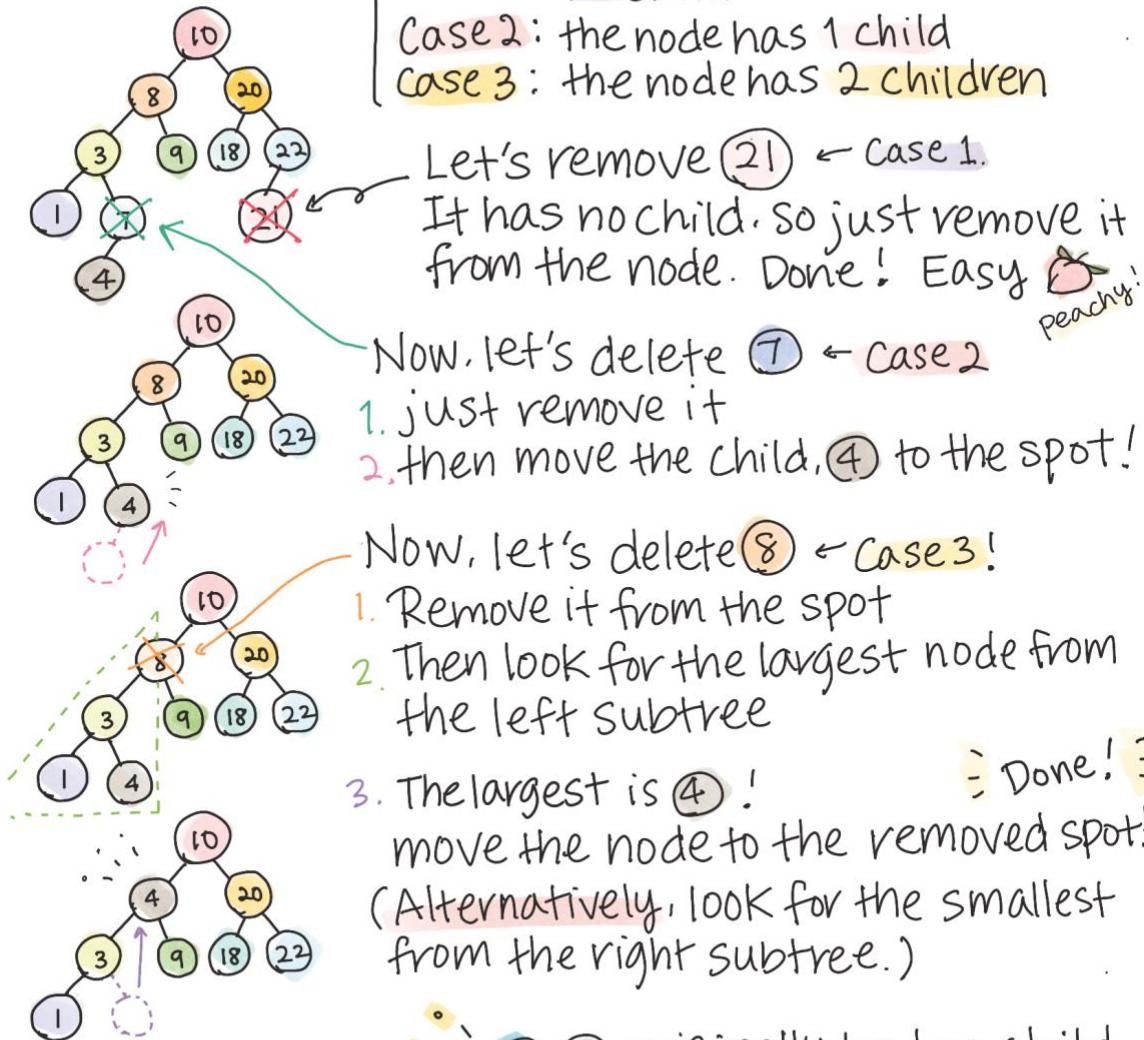


Deletion

Case 1: the to-be-deleted node has no child

Case 2: the node has 1 child

Case 3: the node has 2 children



Complexity:
Ave. $O(\log n)$
Worst. $O(n)$

- ④ originally had no child.
but if it has children?
→ Repeat the process!
 - ↳ Find the largest from left subtree. Move it
 - ↳ Find the largest from left subtree...
- Recursive

Data Structures Hash Table

- A hash table is used to index large amount of data
- Quick key-value lookup. $O(1)$ on average
 - ↳ Faster than brute-force linear search

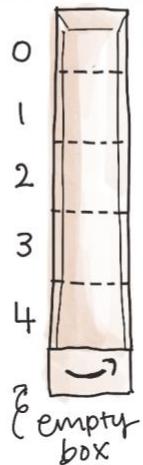
- ① Let's create an array of size 5.

We're going to add data.

Key = "Tabby"

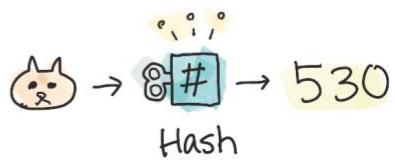
Value = "pizza"

Some data
Let's say, favorite food!



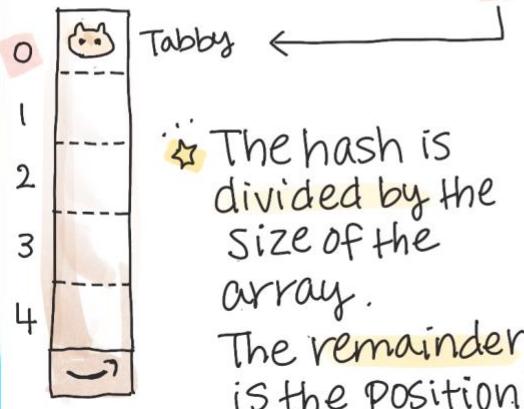
- ② Calculate the hash value by using the key, "Tabby".

e.g. ASCII code, MD5, SHA1



- ③ Use modulo to pick a position in the array!

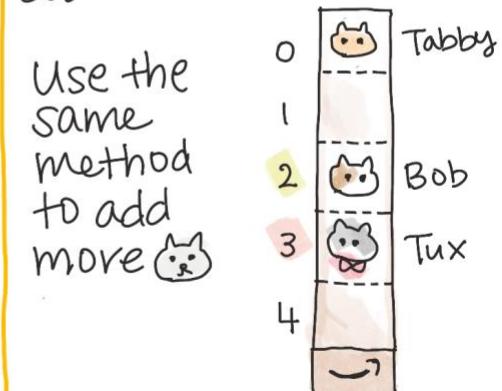
$$\text{Cat} \rightarrow 8 \# \rightarrow 530 \% 5 = 0$$



- ④ Let's add more data.

$$\begin{aligned} \text{Cat} \rightarrow 8 \# &\rightarrow 353 \% 5 = 3 \\ \text{Tux} \rightarrow 8 \# &\rightarrow 307 \% 5 = 2 \\ \text{Bob} \rightarrow 8 \# &\rightarrow 307 \% 5 = 2 \end{aligned}$$

Use the same method to add more .



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★ Collision!

Hash Table

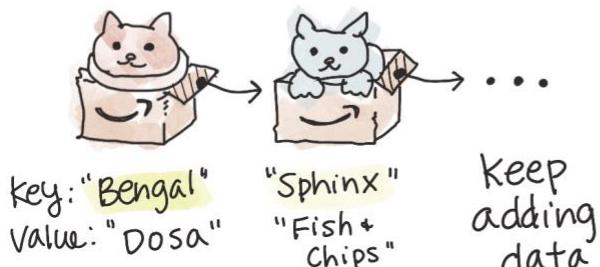
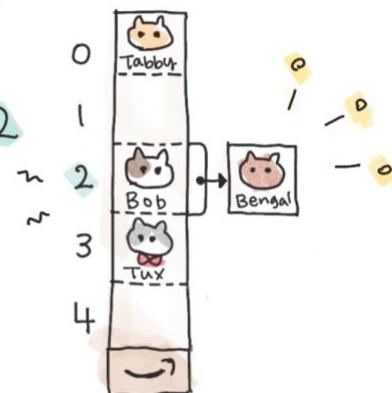
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Now we want to add more data.

Let's add "Bengal".

"Bengal" → 8 # → 617 % 5 = 2

But [2] slot has been taken by "Bob" already! = collision!
so let's chain Bengal next to Bob! = chaining



Searching for data

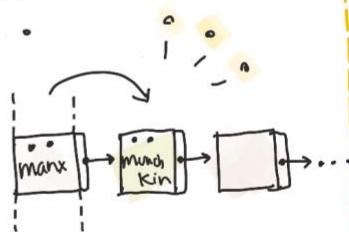
★ Let's look up the value for "Bob"

- ① Get the hash → 307
- ② Get the index → 307 % 5 = 2
- ③ Look up Array [2] → found!

← O(1)

★ Let's look up "munchkin"

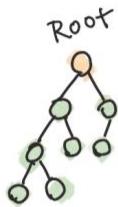
- ① Hash → 861
 - ② Index → 861 % 5 = 1
 - ③ Array [1] → "manx"
 - ④ Operate a linear-search to find "munchkin"
- ↳ Average O(n)



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Data Structures

Binary Heap



Binary tree

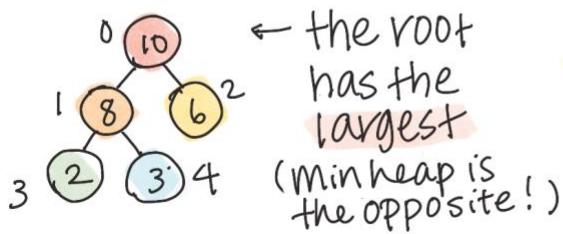
tree data structure
each node has at most 2 children

Binary search tree

Binary heap :

- Complete tree
- Min heap or max heap
- used for priority queue, heap sort etc.

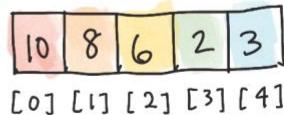
Max heap



the root
has the
largest

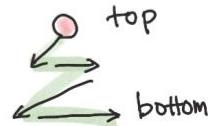
(min heap is
the opposite!)

in array



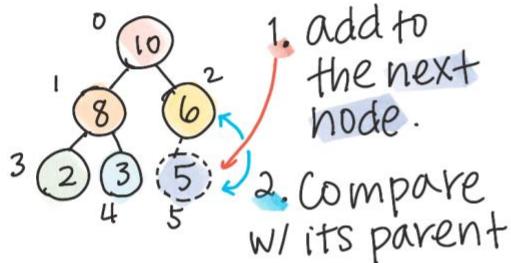
each node has 0 - 2 children

always fill top → bottom, left → right



Insertion

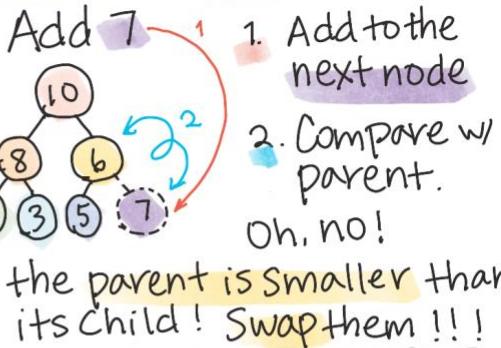
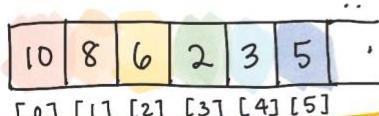
Let's add 5 to the heap!



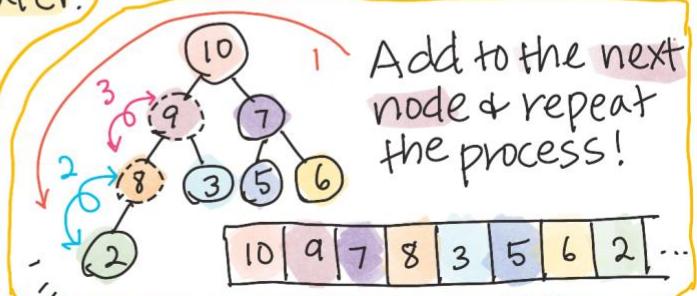
3. the parent is greater.

Cool, it's done!

Let's add more!



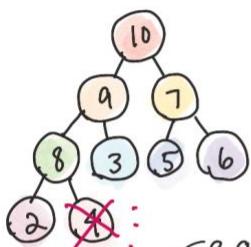
Add to the next node & repeat the process!



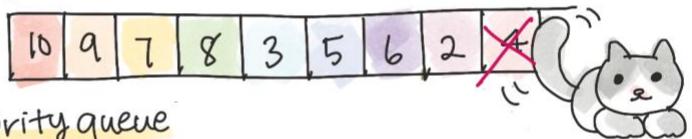
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Data Structures Binary Heap

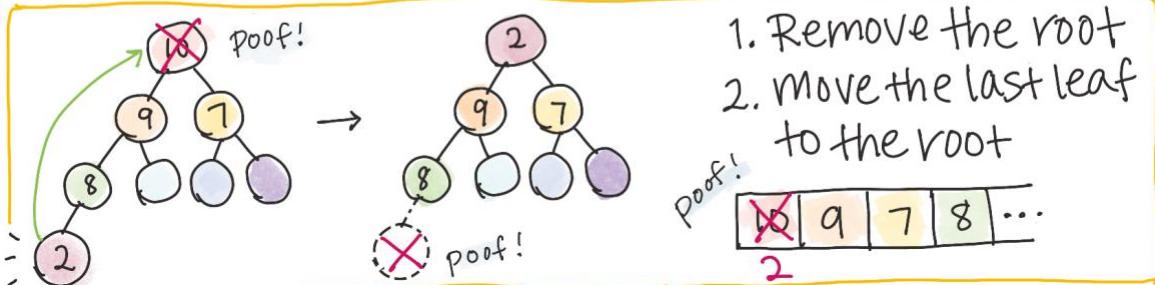
Heap Deletion



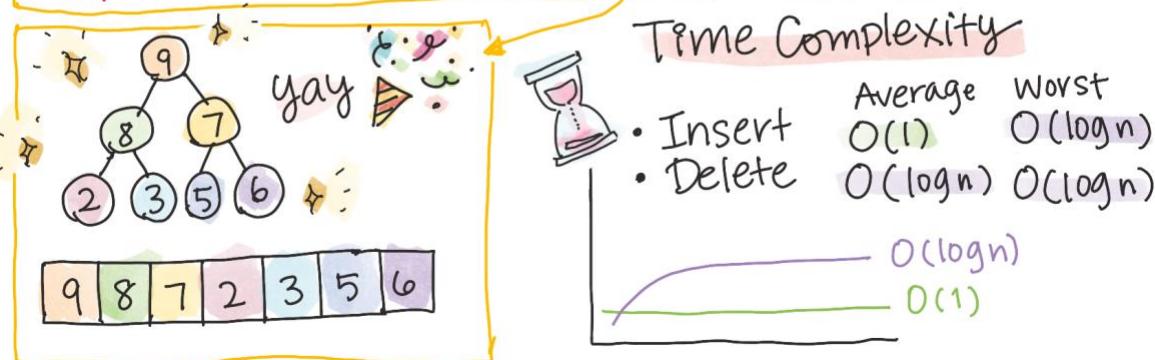
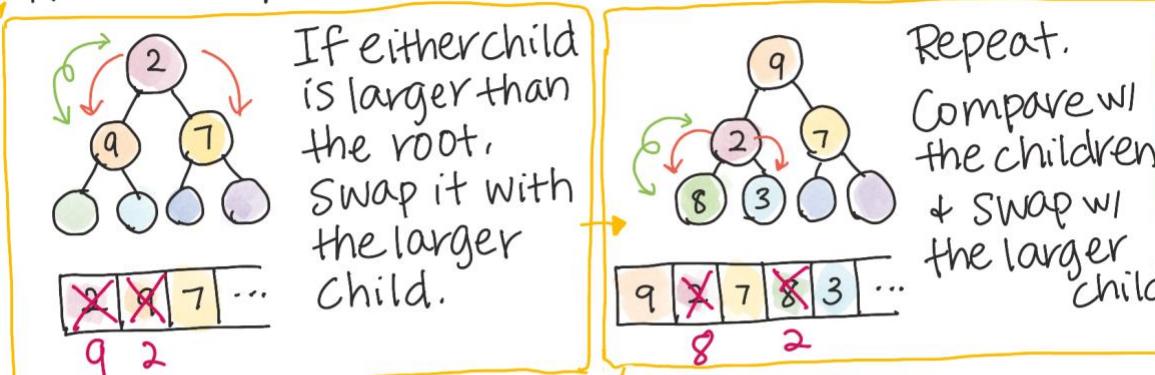
If you want to delete the last leaf,
just delete it!



e.g. priority queue
But typically, you would delete root + heapify!



Now, let's place them in the correct order!



Time Complexity

- Insert $O(1)$
- Delete $O(\log n)$

Average $O(1)$

Worst $O(\log n)$

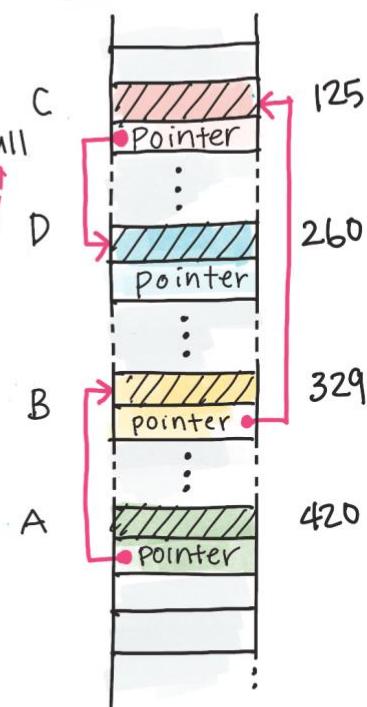
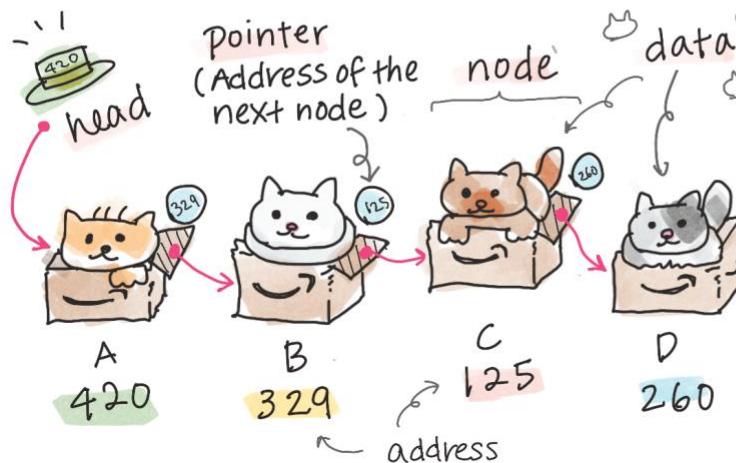
$O(\log n)$

$O(1)$

Linked list

Array & Linked List

- =★ a linear data structure
- =★ each element is a separated object
+ elements are linked w/ pointers



- ★ Unlike an array, Linked List elements are not stored in Contiguous locations.

Yay!	Meh!
♥ Dynamic data	☒ No random access memory.
= Size can grow or shrink	→ Need to traverse n times
♥ Insert & delete element are flexible.	→ time complexity is O(n). array is O(1)
→ no need to shift nodes like array insertion	☒ Reverse traverse is hard
♥ Memory is allocated at runtime	

Data Structures

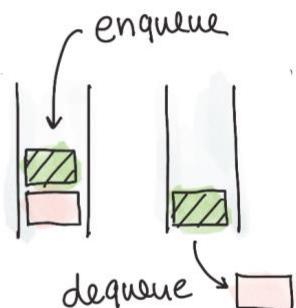
Stack & Queue

LIFO

FIFO

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A queue is a FIFO (First-in-First-out) data structure, where an element added first (= enqueue) gets removed first (= dequeue)



★ Stack as an array in JS

```
let queue = [ ];  
queue.push('Simba'); // ['Simba']  
queue.push('Nyan'); // ['Simba', 'Nyan']  
queue.push('Maru'); // ['Simba', 'Nyan', 'Maru']  
let eater = queue.shift(); // eater is 'Simba'
```

if you
queue.unshift()
(`badKitty`),
then the cat cuts
into the front
of line!

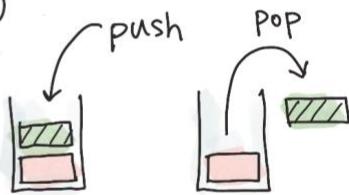
Time Complexity
should be O(1) for
both enqueue + dequeue but JS shift() is slower!

Data Structures

Stack & Queue

LIFO FIFO
@girle_mac

A stack is a LIFO (Last-in-First-out) data structure, where an element added last (=push) gets removed first (=pop)



just like a stack of ice cream scoops!



Stack as an array in



arrays in JavaScript
are dynamic!

```
let Stack = [ ];                                    stack is:  
Stack.push('mint choc'); // ['mint choc']  
Stack.push('vanilla'); // ['mint choc', 'vanilla']  
Stack.push('strawberry'); // ['mint choc', 'vanilla',  
                                                       'strawberry']  
let eaten = Stack.pop(); // eaten is  
                                                       'strawberry'  
                                                       ['mint choc', 'vanilla']
```

Time Complexity is O(1)
for both pop + push.